Certain applications of ultraproducts

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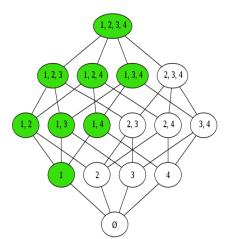
 $3.A \in \mathcal{F} \land A \subseteq B \Rightarrow B \in \mathcal{F}$ (bigger set then a big set is big).

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Lemma

Let U be a filter on I. Then \sim_U is equivalence relation on $\prod A_i$. Conversely, if \sim_U is equivalence relation on $\prod A_i$ then U filter, if $card(A_i) \geqslant 3$, $\forall i \in I$.

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- If R is relation symbol, ar(R) = n, then

$$\langle f_{1_U}, \ldots, f_{n_U} \rangle \in R^{\mathcal{A}} \quad iff \quad \{i \in I \mid \langle f_1(i), \ldots, f_n(i) \rangle \in R^{\mathcal{A}_i}\} \in U;$$

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$$F^{\mathcal{A}}(f_{1_U},\ldots,f_{n_U})=g_U\quad \text{iff}\quad \{i\in I\mid F^{\mathcal{A}_i}(f_1(i),\ldots,f_n(i))=g(i)\}\in U;$$

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• If c is constant, then

$$c^{\mathcal{A}} = \{c^{\mathcal{A}_i} \mid i \in I\}_{\mathcal{U}}.$$

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Ultraproduct construction
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(Łoś's theorem)

For every formula $\varphi(x_1,\ldots,x_n)$ of language \mathcal{L} we have:

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Every colletion of subsets of X with f.i.p. can be extended to ultrafilter on X.

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• infinitesimal:

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• "non-standard large number":

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Let U be an ultrafilter containing \mathcal{J} . Then, $\prod_{P \in S_{\omega}(T)} M_P/U \models T$, because for all

$$\varphi \in T$$

$${P \in S_{\omega}(T) \mid M_P \models \varphi} = J_{{\varphi}} \in U.$$



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Thank you for your attention!