

### Experimental Evaluation of Time-Memory Trade-Off Attack

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## **Encription scenario**



## **Trivium chiper**



- Stream cipher
- Generate keystream
- Uses 80-bit secret key and an 80-bit initial vector (IV)
- Two phases



## Nottation



- cipher C = E(P, K)
- where P is plaintext block of size n
- C is ciphertext block of size n
- K is key of size k







# Inversion of one-way function?

- f(x)=y, if y is known how to find x?
- There are the following two straightforward approaches for recovering the argument given the corresponding image generated by one-way function where the inversion is a hard problem:
  - (i) exhaustive search over all possible arguments;
  - (ii) employing a code-book with all possible argument-image pairs.



## TMTO Attack (Hellman, 1981)

- Compromises the above two extreme approaches
- Precomputation phase: For a given plaintext P:
  - precompute (ideally all) pairs key-ciphertext {K<sub>i</sub>, C<sub>i</sub>};
  - store only some of them in the table.
- Online phase:
  - perform some computations;
  - lookup the table and find the key K.
    - time T = N<sup>2/3</sup>
    - memory  $M = N^{2/3}$

## **Precomputation phase**



- Form a m × t matrix that tries to cover the whole search space which is composed of all the possible comibations of key vector as follows:
  - Step 1. Randomly generate m startpoints of the chains, each point is represented like vector of N bits length.
  - Step 2. Make it the next point in the chain which is the output from Trivium function and update the s register with this point.
  - Step 3. Iterate Step (2) t times on each startpoint respectively.
  - Step 4. Store the pairs of startpoints and endpoints (SPj ,EPj), where j = 1, ...,m in the matrix.





## **TMTO** Issues



- Merging chains
- Cycles in one chain
- Pre-computation is lots of work
- Success is not assured
- There is some repetition rate
- Goal: Try to minimize reppetition rate in precomputing phase, as much as possible.

## **Rainbow tables**



- Several sub-tables
- For each tables generate random vector
- Translating state for coresponding vector
- XOR operation translation
- Dimension of search space is the same like in case of one table



## **Experiments and results**



Number of tables	Repetiton rate [%]*
1	(57.14, 64.43)
2	(54.71, 61.57)
4	(47.14, 58.90)
8	(35.19, 41.62)

\* percent of size of search space

#### **Conclusion:**

The larger the table is, the higher is the probability that a new chain has an intersection with previous chains.



## Thank you for attention!

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