Towards the Verification of Industry 4.0 Applications

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Industry 4.0

New manufacturing paradigm where devices (IoT devices, sensors) are highly interconnected, forming a cyber-physical system.

Smart Factory

- Key features: Interoperability, Information Transparency (raw data into high level data), Decentralized Decisions, Technical Assistance;
- **Examples:** Plug and Produce; using the cloud to optimize production, easy configuration of new production variants;
- Fortiss Future Factory: https:

//www.youtube.com/watch?time_continue=202&v=Tkcv-mbhYqk

4Diac

Programming Industry 4.0 Applications using Models: 4Diac framework.

An application is composed by **function block** (FB) connected by **links**. Links transmit **events** and **data** between FBs.



The **behavior** of the FB is specified by a (**deterministic**) Mealy Machine.

4Diac implements the **IEC 61499** standard for distributed industrial processes.

Application: Simplified PickNPlace





Cyber Attacks can cause Catastrophic Events.

- Cyber-attacks were able to take control of a German mill;
- Ukranian power-station disabled by a cyber-attack;
- Jeep hijacked by hackers;
- Cyber-attacks are able to disable car safety mechanisms.

Levels of Abstraction

- Application Level: Logical level without taking into account the deployment nor the network topology of the system.
- System Level: FBs are deployed into devices. The network topology is abstract. (The colours of the FBs in 4Diac.)
- Network Level: FBs are deployed into devices which are connected through a network topology. (Currently being implemented in 4Diac.)



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Vast number of examples where formal verification can help, including intruder models.

New challenges for formal verification. Devices are weak devices. Powerful Dolev-Yao intruders are not realistic. We need new intruder models with simpler decision problems.

Additionally, one can use

 network defences, such as, filtering suspicious messages.

4Diac Modeling by Example Overview

• We formally specified **an extension with security concerns** of 4Diac semantics in **Maude**.

Features:

- Feature 1 Specification of Applications: Model check for logic errors in the application;
- Feature 2 Deployment: Deployment of applications in systems (sets of devices) as well as in networks (sets of devices connected through switches);
- Feature 3 Intruders: Specification of intruders that can inject messages into channels;
- Feature 4 Defensive Wrappers: Specification of wrappers relying on signed messages;





Bad State: Arm in state mvL and Vac in state off as may lead to the safety hazard of dropping the cap.

Encoding of example in Maude:

Configurations

appInit = [[app, fb(vac, st("off"), none, none) fb(track, st("L"), none, none) fb(ctl, st("init"), none, none) {{ctl,inEv("start")},ev("start")},none]] Input messages to be processed

State Transitions specified as equations

tr(st("init"), st("LOff"), inEv("start") is ev("start"), outEv("GoR") :~ev("GoR"))

Encoding of example in Maude:

Operational Semantics specified as rewrite rules over configurations

crl[app-exe1]: [[id,fbs,iMsgs,none]]
 =>
 [[id,deliverToFBs(fbs,iMsgs),none,none]]
 if not (iMsgs == none) .



May lead to a bad state depending on which event is processed first.

- Feature 1 Specification of Applications: Model check for logic errors in the application;
- badappInit is the configuration with the bad controller.

search badappInit $=>^*$ app such that badState(app).

Finds 4 ways of reaching a bad state in less than 2ms traversing 35 states.

• appInit is the configuration with the good controller.

search appInit $=>^*$ app such that badState(app).

Determines that there is no way to reach a bad state in 1ms traversing 36 states.

 Feature 2 - Deployment: Deployment of applications in systems (sets of devices)...;



Abstract Network

Deployment is computed automatically from the applications and a function block to device mapping.

 Feature 2 - Deployment: Deployment of applications in systems (sets of devices)...;



 Feature 3 - Intruders: Specification of intruders that can inject messages into channels;



- Intruder can inject at any time any anyone of the messages Msg 1, ..., Msg n in the system.
- In the current implementation, intruders can only use a message once. On the one hand, it reduces state space, but on the other hand, it reduces the capabilities of the intruder.

 Feature 3 - Intruders: Specification of intruders that can inject messages into channels;



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Configuration: wsysIntruderNoPol

Intruder possess only one message:

 $\{ dev3, \{ ctl, outEv("VacOff") \} \}, \{ dev1, \{ vac, inEv("VacOff") \} \}, ev("VacOff") \} .$

search wsysIntruderNoPol =>* wsys such that badState(wsys).

Finds 7 attacks in less than 4ms traversing 154 states.

• Feature 4 - Defensive Wrappers: Specifcation of wrappers relying on signed messages;



For example, device with Vac in policy: {(VacOff, Device 1}.

For example, out policy of Device 1 with the controller: {VacOff}.

• Feature 4 - Defensive Wrappers: Specifcation of wrappers relying on signed messages;

Configuration: wsysIntruderPol

- Intruder possess only one message: {{dev3,{ctl,outEv("VacOff")}},{dev1,{vac,inEv("VacOff")}},ev("VacOff")}.
- Devices with defensive wrappers.

search wsysIntruderPol =>* wsys such that badState(wsys).

Check that there are no attack in less than 4ms traversing 38 states.

Feature 2 - Deployment: ... as well as in networks (sets of devices connected through switches);



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Configuration: netIntruderNoPol

- Intruder possess only one message: {{dev3,{ctl,outEv("VacOff")}},{dev1,{vac,inEv("VacOff")}},ev("VacOff")}.
- Devices without defensive wrappers;
- Devices connected through one switch.
 search netIntruderNoPol =>* nsys such that badState(nsys) .
 Finds 12 attacks in around 60ms traversing 488 states.

Configuration: netIntruderPol

• Devices with defensive wrappers;

search netIntruderPol $=>^*$ nsys such that badState(nsys). Concludes that there are no attacks in 25ms traversing 200 states.

• Feature 2 - Deployment: ... as well as in networks (sets of devices connected through switches);

Configuration: netIntruderPol

- Intruder possess only one message: {{dev3,{ctl,outEv("VacOff")}},{dev1,{vac,inEv("VacOff")}},ev("VacOff")}.
- Devices with defensive wrappers;
- Devices connected through one switch.

search netIntruderPoI =>* nsys such that badState(nsys).

Checks that no attack is possible in 13ms traversing 98 states.

Other uses of Formal Verification

Deployment Optimization: Encryption is **heavy and undersirable** as it leads to delays. Unfeasible for applications such as the **Tomato Separator**:

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https://www.youtube.com/watch?v=EBddJjYNp0g
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Our Idea (still under development)



Further Questions

Issues / Questions

• Performance:

TSN communication;

What is the impact of signature? Weaker signature schemes?

• Defenses using Network:

Can we use firewall to mitigate attacks, by, e.g., isolate sub-nets? Port Security?

• Intruder Model:

Is it running on a good device? Is it running on its own device? What information can we rely on the HW/wiring to provide that is not fungible?

Future Work

- Model defenses for switches: firewalls?, SDN rules?
- Decision problems and complexity results;
- Include data channels;
- Stronger Intruder Models: Symbolic Analysis
- Stronger defensive wrappers;
- Investigate scalability;
- Extend model with TSN.

Questions?