Probabilistic Reasoning about Typed Combinatory Terms

Simona Kašterović

Faculty of Technical Sciences, University of Novi Sad

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1. Logic of Combinatory Logic (LCL)
   joint work with Silvia Ghilezan

2. Probabilistic Reasoning about Typed Combinatory Terms (PCL)
   joint work with Silvia Ghilezan, Jelena Ivetić, Zoran Ognjanović, Nenad Savić

Outline

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LCL Syntax

- Logic of combinatory logic (*LCL*) - classical propositional logic over simply typed combinatory logic

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$$M : \sigma$$

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Syntax of *LCL*

$$\alpha, \beta := M : \sigma \mid \neg\alpha \mid \alpha \Rightarrow \beta$$

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(Ax 2)	$\mathsf{K} : \sigma \to (\tau \to \sigma)$
(Ax 3)	$\mathsf{I} : \sigma \to \sigma$
(Ax 4)	$(M : \sigma \to \tau) \Rightarrow ((N : \sigma) \Rightarrow (MN : \tau))$
(Ax 5)	$M : \sigma \Rightarrow N : \sigma$, if $M = N$
(Ax 6)	$\alpha \Rightarrow (\beta \Rightarrow \alpha)$
(Ax 7)	$(\alpha \Rightarrow (\beta \Rightarrow \gamma)) \Rightarrow ((\alpha \Rightarrow \beta) \Rightarrow (\alpha \Rightarrow \gamma))$
(Ax 8)	$(\neg\alpha \Rightarrow \neg\beta) \Rightarrow ((\neg\alpha \Rightarrow \beta) \Rightarrow \neg\alpha)$

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- Inference rule:

$$\frac{\alpha \Rightarrow \beta \qquad \alpha}{\beta}\ (\mathsf{MP})$$

LCL Semantics

Applicative structure for *LCL*

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A *LCL*-model is a tuple $\mathcal{M}_\rho = \langle \mathcal{M}, \rho \rangle$, where $\mathcal{M}$ is an applicative structure and ρ is an environment for $\mathcal{M}$.

Soundness and completeness of the axiomatiation for *LCL*

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Completeness of Ax

If $T \models \alpha$, then $T \vdash \alpha$.

Simona Kašterović and Silvia Ghilezan. Logic of combinatory logic, February 2022. Submitted for publication

In addition...

The proposed semantics is a new semantics proven to be sound and complete for combinatory logic.

If $\Gamma \vdash_{\mathsf{CL}} M : \sigma$, then $\Gamma \models M : \sigma$.

Let Γ be a basis. If $\Gamma \models M : \sigma$, then $\Gamma \vdash_{\mathsf{CL}^{=}} M : \sigma$.

LCL is a conservative extension of the simply typed combinatory logic.

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PCL

- *PCL* - a probabilistic system for simply typed combinatory terms;
- Probabilistic logic over *LCL*;
- Probabilistic logic LPP_2

$$\phi := P_{\geq s}\alpha \mid \neg\phi \mid \phi \wedge \phi,$$

  where α is a formula of classical propositional logic and $s \in [0,1] \cap \mathbb{Q}$.
- *PCL* syntax

$$\phi := P_{\geq s}\alpha \mid \neg\phi \mid \phi \wedge \phi,$$

  where α is an *LCL*-formula and $s \in [0,1] \cap \mathbb{Q}$.
- $P_{\geq s}\alpha$ has a meaning "probability that α is true is greater than or equal to s"

Zoran Ognjanović, Miodrag Rašković, and Zoran Marković. Probability Logics - Probability-Based Formalization of Uncertain Reasoning. Springer, 2016.

PCL Axiomatization

LCL axiomatization + axiomatic system for probability logic

Axiomatic system for probability logic:

(1) all instances of the classical propositional tautologies, (atoms are any *PCL*-formulas),

(2) $P_{\geq 0}\alpha$,

(3) $P_{\leq r}\alpha \Rightarrow P_{<s}\alpha$, $s > r$,

(4) $P_{<s}\alpha \Rightarrow P_{\leq s}\alpha$,

(5) $(P_{\geq r}\alpha \wedge P_{\geq s}\beta \wedge P_{\geq 1}(\neg\alpha \vee \neg\beta)) \Rightarrow P_{\geq min\{1,r+s\}}(\alpha \vee \beta)$,

(6) $(P_{\leq r}\alpha \wedge P_{<s}\beta) \Rightarrow P_{<r+s}(\alpha \vee \beta)$, $r + s \leq 1$,

(7) $P_{\geq 1}(\alpha \Rightarrow \beta) \Rightarrow (P_{\geq s}\alpha \Rightarrow P_{\geq s}\beta)$.

PCL Semantics - an idea

$$\mathcal{M} = \langle W, \{\rho_w\}_w, H, \mu \rangle$$

- W is a nonempty set of worlds, where each world is one LCL-applicative structure;
- $\rho_w : V \times \{w\} \to D_w$;
- H is an algebra of subsets of W;
- μ is a finitely additive probability measure defined on H.

S. Ghilezan, J. Ivetić, S. Kašterovć, Z. Ognjanović, and N. Savić. Towards probabilistic reasoning in type theory - the intersection type case. 11th International Symposium, FoIKS 2020, Dortmund, Germany, February 17-21, 2020, Proceedings, volume 12012 of Lecture Notes in Computer Science, pages 122–139. Springer, 2020.

S. Ghilezan, J. Ivetić, S. Kašterović, Z. Ognjanović, and N. Savić. Probabilistic reasoning about typed lambda terms. International Symposium, LFCS 2018, Deerfield Beach, FL, USA, January 8-11, 2018, Proceedings, volume 10703 of Lecture Notes in Computer Science, pages 170–189. Springer, 2018

Ongoing

- Soundness;
- Completeness.

Thank you for your attention!